

Heuristic Algorithms

Master's Degree in Computer Science/Mathematics

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Schedule: **Thursday 14.30 - 16.30 in classroom 503**

Friday 14.30 - 16.30 in classroom 503

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Web page: **<https://homes.di.unimi.it/cordone/courses/2024-ae/2024-ae.html>**

Ariel site: **<https://myariel.unimi.it/course/view.php?id=4466>**

Aims of the course

This course aims to

- 1 show that **heuristic algorithms are not recipes for specific problems:**
heuristics and problems can be matched freely
(of course, with different performance)
- 2 discuss the **common and general aspects** of these algorithms
- 3 teach **how to design** a heuristic for a specific problem
- 4 teach **how to evaluate its performance**

eurisko = I find

It is a word derived from Greek

- inspired by the famous story of Archimedes and the golden crown



but it was

- never used by the ancient Greeks
- coined during the 19th century

Some historical facts

- 4th century CE: Pappus of Alexandria discusses the *analytomenos* (*treasure of analysis*), that is **how to build a mathematical proof**
 - how to move from the hypotheses to the thesis of a theorem
 - how to move from the data to the solution of a geometrical problem
- 17th century: Descartes, Leibnitz *et al.* discuss the *ars inveniendi* (*art of finding*), *i. e.* the **attainment of truth through mathematics**
- 19th century: Bernard Bolzano discusses in detail the most common strategies to **build mathematical proofs** (*Erfindungskunst*)
- 19th-20th century: philosophers, psychologists and economists define **heuristics** as practical and simple **decision rules** that do not aim at an optimal result, but at a **satisficing** one (Simon, 1957)
- 1945: the short essay *How to solve it* by György Pólya comes back to the mathematical meaning of heuristic as an **informal process that leads to prove a thesis or to find a solution**

So, what about *heuristic algorithms*?

Algorithms and heuristics

Some scientific sectors use the two words as opposites:

- **algorithm** as a **formal, deterministic procedure, consisting of a finite sequence of elementary steps**
- **heuristic** as an **informal, creative, open rule**

One could even say that

- **an algorithm is a correctness proof**
- **a heuristic is a bunch of common sense arguments**

In fact, an algorithm has a correctness proof, a heuristic has none

The phrase **heuristic algorithm** is an oxymoron, in some respects

Then what does it mean?

Heuristic algorithms

A heuristic algorithm is an algorithm which does not guarantee a correct solution

Then it is useless!

Quite to the contrary, it can be useful, provided that

- 1 it “costs” much less than a correct algorithm:
this requires a definition of **computational cost** of an algorithm
 - time
 - space
- 2 it “frequently” yields something “close” to the correct solution:
this requires to define a solution space endowed with
 - a **metric** to express a “satisfactory distance” from the correct solution
 - a **probabilistic distribution** to express the “satisfactory frequency” of solutions at a satisfactory distance from the correct solutions

Proofs and algorithms

Mathematical proofs and algorithms are strictly related

- every algorithm has/is a correctness proof
- both are mechanical symbolic transformations from a starting point (hypotheses/data) to an ending point (thesis/solution)
- Turing's undecidability proof mirrors Gödel's incompleteness proof

Heuristics are the construction of both proofs and algorithms

- in case of success, the heuristic is abandoned and the proof preserved
- otherwise, a good heuristic frequently provides a good result, instead of always providing a perfect one

This is the motivation for heuristic algorithms

The focus of this course

The course focuses on heuristic algorithms

- that apply to **Combinatorial Optimization** problems
- that are **solution-based** (as opposed to **model-based**)

So, we limit

- ① the kind of problem
- ② the kind of algorithm

It is still a pretty wide field

Let us further discuss the two limitations

Problem classification

A problem is a question on a mathematical system

Problems can be classified based on the nature of their solution:

- **decision problems**: their solution is either *True* or *False*
- **search problems**: their solution is *any feasible subsystem* (that is, satisfying certain conditions)
- **optimization problems**: their solution is the *minimum or maximum value of an objective function defined on the feasible subsystems*
- **counting problems**: their solution is the *number of feasible subsystems*
- **enumeration problem**: their solution is the *collection of all feasible subsystems*
- ...

We address the combination of optimization and search, that is, we look for the optimal value and a subsystem assuming that value

Optimization/search problems

An optimization/search problem can be represented as

$$\begin{aligned} &\text{opt } f(x) \\ &x \in X \end{aligned}$$

where

- a **solution** x describes **each subsystem** of the problem
- the **feasible region** X (**feasible solution space**) is the **set of subsystems which satisfy given conditions**
- the **objective function** $f : X \rightarrow \mathbb{R}$ **quantitatively measures the quality of each subsystem** ($\text{opt} \in \{\text{min}, \text{max}\}$)

The problem consists in determining

- optimization: the **optimal value** f^* of the objective function:

$$f^* = \underset{x \in X}{\text{opt}} f(x)$$

- search: at least one **optimal solution**, that is a subsystem

$$x^* \in X^* = \underset{x \in X}{\text{arg opt}} f(x) = \left\{ x^* \in X : f(x^*) = \underset{x \in X}{\text{opt}} f(x) \right\}$$

Why optimization/search problem?

Several application fields **require objects or structures** characterized by **very high or very low values** of a suitable **evaluation function**

- *bioinformatics*: the most effective drugs bond with proteins in configurations of **minimal potential energy**
- *social networks*: the best target for a campaign are the **most influential**, **most influential** and **most uncorrelated** groups of individuals
- *machine learning*: the most effective classification systems generate the **simplest** classifications and the **minimum amount of violations**
- *hardware design*: the best logical circuits require the **minimum space** and yield the **minimum delay**
- *parameter estimation*: the best physical models are the ones which reproduce the observations with the **minimum error**
- *finance*: the most effective portfolio management algorithms reproduce the target time series in the **most precise** way

Exact optimality is costly, not always required, or even desirable
(*many heuristic solutions could be preferable to a single exact one*)

Why optimization/search problem?

Other problems can often be reduced to optimization/search problems

- **hard search problems** can be reduced by
 - relaxing the conditions to satisfy, so as to enlarge the feasible region from X to $X' \supset X$ and obtain an easy search problem;
 - introducing a function $d(x)$ to **quantify the distance** of each $x \in X'$ from X ;
 - minimizing $d(x)$ to find x^* such that $d(x^*) = 0 \Leftrightarrow x^* \in X$.
- some **decision problems** concern the existence of feasible subsystems, and are equivalent to search problems
(*finding the subsystem proves that it exists*)
- some **enumeration problems** concern the search for subsystems with “good” values of conflicting objective functions (Pareto frontier) and allow **direct adaptations** of optimization/search algorithms

Such reductions are often possible and useful, though not always

Combinatorial Optimization (CO)

A problem is a CO problem when **the feasible region X is a finite set**, that is, **it has a finite number of feasible solutions**

This looks like a very restrictive assumption

However, the study of CO problems can be useful more in general:

- ① infinite discrete problems can have a finite set of interesting solutions
- ② some continuous problems can be reduced to CO problems
(e. g., *Linear Programming, Maximum Flow, Minimum Cost Flow*)
- ③ continuous problems can be reduced to discrete ones by sampling
(*usually not very effective*)
- ④ ideas conceived for CO problems can be extended to other problems
(*often quite effective*)

Model-based heuristics

They describe the feasible region X with a “model”

A typical example is a Mathematical Programming formulation

$$\begin{array}{l} \text{opt } f(x) \\ x \in X \end{array} \quad \longrightarrow \quad \begin{array}{l} \min \phi(\xi) \\ g_i(\xi) \leq 0 \quad i = 1, \dots, m \end{array}$$

where

- $\xi \in \mathbb{R}^n$, that is, a solution is a vector of n real values
- $X = \{\xi \in \mathbb{R}^n : g_i(\xi) \leq 0, i = 1, \dots, m\}$, that is, the feasible region is the set of vectors which satisfy all the inequalities (constraints)

Model-based heuristics exploit the information derived from the model, that is the analytical properties of functions ϕ and g_i ($i = 1, \dots, m$)

Other models can be based on SAT, etc. . .

We will not use these tools

An alternative definition of CO

A problem is a CO problem when:

- ① the number of feasible solutions is finite
- ② the feasible region is $X \subseteq 2^B$ for a given finite ground set B , that is, the feasible solutions are all subsets of the ground set that satisfy suitable conditions

The two definitions are equivalent:

$2 \Rightarrow 1$: if the ground set B is finite, every collection $X \subseteq 2^B$ is finite

$1 \Rightarrow 2$: if the number of feasible solutions is finite, define B as their set and the feasible region X as the collection of all singletons of B
(a “solution” is a set containing a single solution)

In general, the sophisticated definition allows a deeper analysis, because

- X is not simply enumerated
- X is defined in a compact and significant way

Solution-based heuristics: a classification for CO problems

Solution-based heuristics consider solutions as subsets of the ground set

① **constructive/destructive heuristics:**

- they start from an extremely simple subset (respectively, \emptyset or B)
- they add/remove elements until they obtain the desired solution

② **exchange heuristics:**

- they start from a subset obtained in any way
- they exchange elements until they obtain the desired solution

③ **recombination heuristics:**

- they start from a population of subsets obtained in any way
- they recombine different subsets producing a new population

Heuristic designers can creatively combine elements from different classes

Randomization and memory

Two other distinctions concern

- the use of **randomization**:
 - **deterministic** heuristics, whose input includes only certain information
 - **randomized** heuristics, whose input includes **pseudorandom numbers**
(*they are deterministic algorithms anyway*)
- the use of **memory**:
 - heuristics whose input includes only current information
 - heuristics whose input also includes **previously generated solutions**

These distinctions are independent from the previous classification

Metaheuristics (from the Greek, “beyond heuristics”) is the common name for **heuristic algorithms with randomization and/or memory**

Risks to beware of

- 1 **reverential or trendy attitude**, that is choosing an algorithm based on the social context, instead of the problem
- 2 **magic attitude**, that is trusting a method on the basis of an analogy with physical and natural phenomena
- 3 **heuristic integralism**, that is using a heuristic for a problem which admits exact algorithms
- 4 **number crunching**, that is performing sophisticated and complex computations with unreliable numbers
- 5 **SUV attitude**, that is relying on hardware power
- 6 **overcomplication**, that is introducing redundant components and parameters, as if that could only improve the result
- 7 **overfitting**, that is adapting components and parameters of the algorithm to the specific dataset used in the experimental evaluation

It is fundamental to

- free oneself from prejudices
- evaluate the performance of the algorithm in a scientific way
- distinguish the contribution of each component of the algorithm
- efficiently implement each component of the algorithm

The *Analytics and Optimization* track

This course belongs to the *Analytics and Optimization* track:

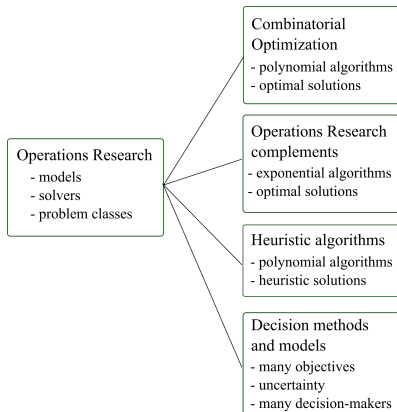
- this gives a specific slant to the presentation of the subject
- but the course can be easily attended by students of other tracks

The focus of the track is on practical decisions where

- a large amount of data must be taken into account
- the possible choices are many
- the costs of a wrong choice are high

The correct strategy is **first make a model, then compute, finally decide** and nowadays this is supported by

- **Big Data**: huge amounts of precise, structured and cheap data, from which to extract information
- **Cloud Computing**: pervasive capacity to access and process data
- **Business Analytics**: a business culture open to the use of models
- **new theory**: online, stochastic, robust programming, etc. . .



But we have already seen that

- heuristic algorithms are useful for any application
- we do not require models

The course is open to any student with these prerequisites

- C programming (laboratory)
- Algorithms and data structures (preferential)