

Foundations of Operations Research

Master of Science in Computer Engineering

Roberto Cordone

roberto.cordone@unimi.it

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Thursday 10.15 - 13.15

<http://homes.di.unimi.it/~cordone/courses/2014-for/2014-for.html>



A sample problem

$$\begin{aligned}\min f &= 5x_1 + 4x_2 + 3x_3 \\ 2x_1 + 3x_2 + x_3 + x_4 &= 5 \\ 4x_1 + x_2 + 2x_3 + x_5 &= 11 \\ 3x_1 + 4x_2 + 2x_3 + x_6 &= 8 \\ x_1, x_2, x_3, x_4, x_5, x_6 &\geq 0\end{aligned}$$

There are 20 basic solutions, but it is easy to find the optimal one

- 1 matrix A shows an obvious basic submatrix $B = I$ (columns 4, 5 and 6) and the basic cost subvector is null: $c_B = 0$:
 - the corresponding basic solution is $\bar{x}_N = 0 \Rightarrow x_B = B^{-1}b = b$
 - the value of the objective function is $f(\bar{x}) = c_B \bar{x}_B + c_N \bar{x}_N + d = d$
- 2 the right-hand-side vector b is nonnegative: the obvious basic solution is feasible
- 3 the nonbasic cost subvector c_N is nonnegative: $c_N \geq 0$
a generic feasible solution x costs $f(x) = c_N^T x_N + d \geq d$

Therefore, the obvious basic solution $\bar{x} = \begin{bmatrix} b \\ 0 \end{bmatrix}$ is optimal

Not all LP problems enjoy such properties, but all feasible and bounded LP problems admit an equivalent problem that enjoys them

Basic canonical form

A **basic canonical form** is a special standard form in which

- the coefficient matrix A includes a submatrix $B = I$
- the cost vector c includes a subvector $c_B = 0$
- B and c_B correspond to the same variables

This is the first of the three conditions mentioned above

An LP problem admits a basic canonical form for each of its bases

A direct way to obtain it is to multiply the equality constraints by B^{-1}

$$B^{-1} A x = B^{-1} [B \mid N] \begin{bmatrix} x_B \\ x_N \end{bmatrix} = B^{-1} b \Rightarrow x_B + B^{-1} N x_N = B^{-1} b$$

and then replace x_B in the expression of the objective function

$$f(x) = c^T x = [c_B^T \mid c_N^T] \begin{bmatrix} x_B \\ x_N \end{bmatrix} = (c_N^T - c_B^T B^{-1} N) x_N + c_B^T B^{-1} b + d$$

From the basic canonical form, it is easy to obtain

- the associated basic solution: $\bar{x} = \bar{b} = \begin{bmatrix} B^{-1} b \\ 0 \end{bmatrix}$
- the associated objective value: $f(\bar{x}) = \bar{d} = c_B^T B^{-1} b + d$
- the associated reduced cost vector: $\bar{c}^T = [0 \mid c_N^T - c_B^T B^{-1} N]$

An algorithmic idea

Given an LP problem in standard form

- 1 put it into a basic canonical form (*How? Premultiplying by B^{-1} ?*)
- 2 if $\bar{b} \geq 0$, the associated basic solution \bar{x} is feasible;
otherwise, put the problem into another basic canonical form
and go back to step 1 (*How? What form? Is it always possible?*)
- 3 if $\bar{c}_N \geq 0$, the associated feasible basic solution \bar{x} is also optimal;
otherwise, put the problem into another basic canonical form
and go back to step 2 (*How? What form? Is it always possible?*)

The simplex method examines a sequence of basic canonical forms (solutions) in two phases

- 1 achieving a basic feasible solution by subsequent reductions of the unfeasibility (suitable conditions reveal whether the problem is unfeasible)
- 2 achieving a basic optimal solution by subsequent reductions of the objective (suitable conditions reveal whether the problem is unbounded)

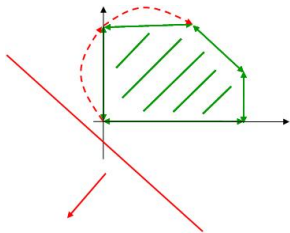
The simplex method: general idea

The idea is to **move from a basic canonical form to an “adjacent” one**

- ① a single nonbasic column gets into the basis; correspondingly, a single nonbasic variable changes from zero to nonnegative
- ② a single basic column gets out of the basis; correspondingly, a single basic variable changes from nonnegative to zero

From a geometrical point of view, the idea is that

the current solution moves from a vertex to an “adjacent” vertex



The simplex method consists of two phases:

- ① the current basic solution gets closer to the feasible region
- ② the basic solution keeps feasible, and the objective function improves

Obtaining a basic canonical form

The following problem is in standard form, but not in any basic canonical form

$$\begin{aligned}\min f &= -6x_1 - 5x_3 - \frac{7}{3}x_4 + \frac{8}{3}x_5 + \frac{26}{3} \\ 3x_1 + 3x_3 + x_4 - 2x_5 &= 2 \\ 3x_1 + 3x_2 + x_4 + x_5 &= 8 \\ 6x_3 + x_4 - 5x_5 + 3x_6 &= 1 \\ x_1, \dots, x_6 &\geq 0\end{aligned}$$

Arbitrarily selecting columns (1, 2, 6), one obtains a basis B :

$$B = \begin{bmatrix} 3 & 0 & 0 \\ 3 & 3 & 0 \\ 0 & 0 & 3 \end{bmatrix} \Rightarrow B^{-1} = \frac{1}{27} \cdot \begin{bmatrix} 9 & 0 & 0 \\ -9 & 9 & 0 \\ 0 & 0 & 9 \end{bmatrix}$$

The corresponding canonical form can be easily built by simple matrix operations

$$\bar{A} = B^{-1} \cdot A = \begin{bmatrix} 1 & 0 & 1 & 1/3 & 2/3 & 0 \\ 0 & 1 & -1 & 0 & 1 & 0 \\ 0 & 0 & 2 & 1/3 & -5/3 & 1 \end{bmatrix} \quad \bar{b} = B^{-1}b = \begin{bmatrix} 2/3 \\ 2 \\ 1/3 \end{bmatrix}$$

The corresponding reduced cost vector and constant term are

$$\bar{c}^T = [0 \quad 0 \quad 1 \quad -1/3 \quad -4/3 \quad 0] \quad \bar{d} = c_B^T B^{-1}b + d = \frac{14}{3}$$

However, there is a simpler and more efficient way to proceed

Equivalent equalities

Any LP problem in standard form can be transformed into an equivalent one by

- 1 multiplying both sides of an equality by a constant coefficient

$$3x_1 + 3x_3 + x_4 - 2x_5 = 2 \Leftrightarrow 1x_1 + 1x_3 + \frac{1}{3}x_4 - \frac{2}{3}x_5 = \frac{2}{3}$$

- 2 using an equality to determine a variable in terms of the other ones and replacing the variable everywhere else

$$\min f = -6x_1 - 5x_3 - \frac{7}{3}x_4 + \frac{8}{3}x_5 + \frac{26}{3}$$

$$3x_1 + 3x_3 + x_4 - 2x_5 = 2$$

$$3x_1 + 3x_2 + x_4 + x_5 = 8$$

$$6x_3 + x_4 - 5x_5 + 3x_6 = 1$$

$$x_1, \dots, x_6 \geq 0$$

Use the first constraint to determine $x_1 = \frac{2 - 3x_3 - x_4 + 2x_5}{3}$

$$\min f = x_3 - \frac{1}{3}x_4 - \frac{4}{3}x_5 + \frac{14}{3}$$

$$x_1 + x_3 + \frac{1}{3}x_4 - \frac{2}{3}x_5 = \frac{2}{3}$$

$$3x_2 - 3x_3 + 3x_5 = 6$$

$$6x_3 + x_4 - 5x_5 + 3x_6 = 1$$

$$x_1, \dots, x_6 \geq 0$$

Equivalent equalities

The same transformation can be done

- 1 multiplying both sides of an equality by a constant coefficient

$$3x_1 + 3x_3 + x_4 - 2x_5 = 2 \Leftrightarrow 1x_1 + 1x_3 + \frac{1}{3}x_4 - \frac{2}{3}x_5 = \frac{2}{3}$$

- 2 subtracting from an equality another equality multiplied by a suitable constant

$$\begin{aligned} \min f &= -6x_1 - 5x_3 - \frac{7}{3}x_4 + \frac{8}{3}x_5 &+ \frac{26}{3} \\ 3x_1 + 3x_3 + x_4 - 2x_5 &= 2 \\ 3x_1 + 3x_2 + x_4 + x_5 &= 8 \\ 6x_3 + x_4 - 5x_5 + 3x_6 &= 1 \\ x_1, \dots, x_6 &\geq 0 \end{aligned}$$

Multiply the first constraint by 1/3
Subtract the first constraint

- multiplied by 3
from the second constraint
- multiplied by 0
(x_1 is already absent)
from the third constraint

$$\begin{aligned} \min f &= x_3 - \frac{1}{3}x_4 - \frac{4}{3}x_5 &+ \frac{14}{3} \\ x_1 + x_3 + \frac{1}{3}x_4 - \frac{2}{3}x_5 &= \frac{2}{3} \\ 3x_2 - 3x_3 + 3x_5 &= 6 \\ 6x_3 + x_4 - 5x_5 + 3x_6 &= 1 \\ x_1, \dots, x_6 &\geq 0 \end{aligned}$$

Consider the objective function
as an equality: $f - d = c^T x$ and
subtract from it the first constraint
multiplied by -6

The *tableau* representation

A common representation of an *LP* problem in standard form is the so called *tableau*, that is a matrix with $m + 1$ rows and $n + 1$ columns

- row 0 corresponds to the objective function: $a_{0j} := c_j$ ($j = 1, \dots, n$)
rows $1, \dots, m$ correspond to the constraints
- column 0 corresponds to the right-hand sides: $a_{i0} := b_i$ ($i = 1, \dots, m$)
columns $1, \dots, n$ correspond to the variables

Since $c^T x = f - d$, we set $a_{00} := -d$

$$\begin{aligned} \min f &= c^T x + d \\ Ax &= b \\ x &\geq 0 \end{aligned}$$

$-d$	c^T
b	A

$$\begin{aligned} \min f &= -6x_1 - 5x_3 - \frac{7}{3}x_4 + \frac{8}{3}x_5 + \frac{26}{3} \\ 3x_1 + 3x_3 + x_4 - 2x_5 &= 2 \\ 3x_1 + 3x_2 + x_4 + x_5 &= 8 \\ 6x_3 + x_4 - 5x_5 + 3x_6 &= 1 \\ x_1, \dots, x_6 &\geq 0 \end{aligned}$$

$-\frac{26}{3}$	-6	0	-5	$-\frac{7}{3}$	$\frac{8}{3}$	0
2	3	0	3	1	-2	0
8	3	3	0	1	1	0
1	0	0	6	1	-5	3

The *pivot* operation

The *pivot operation* represents on the *tableau* the substitution of a variable
It is based on

- a *pivot column* j : the variable x_j expressed in terms of the others
- a *pivot row* i : the constraint (a_i^T, b_i) used to determine the variable

The *pivot* is the element a_{ij} identified by the *pivot* row and column

The *pivot* operation consists in:

- dividing the *pivot* row $[b_i \mid a_i^T]$ by the *pivot* element a_{ij} ($a_{ij} \neq 0!$)

$$\bar{a}_i^T := \frac{a_i^T}{a_{ij}} \quad (\text{note: } b_i = a_{i0} \Rightarrow \bar{b}_i := \frac{b_i}{a_{ij}})$$

- subtracting the *pivot* row multiplied by a suitable value from all other rows

$$\bar{a}_h := a_h - \frac{a_{hj}}{a_{ij}} a_i \quad \forall h \neq i$$

$$(\text{note: } c_j = a_{0j} \Rightarrow \bar{c} := c - \frac{c_j}{a_{ij}} a_i \text{ and } d = a_{00} \Rightarrow \bar{d} := d - \frac{c_j}{a_{ij}} b_i)$$

For the *pivot* column j , the operation yields $\bar{c}_j = \bar{a}_{0j} := 0$ and $\bar{a}_{hj} := 0$

The effect of a *pivot* operation is that the *pivot column* gets into the basis

Obtaining a basic canonical form

Given an LP problem and a basis, **one obtains the associated basic canonical form applying the transformation to each column of the basis**; for each basic variable:

- the coefficient in the objective function becomes $\bar{c}_j := 0$
- the *pivot* element becomes $\bar{a}_{ij} := 1$
- the coefficients of the *pivot* column in the nonpivot rows become $\bar{a}_{hj} := 0, \forall h \neq i$

$\Rightarrow \bar{B} = I$ and $\bar{c}_B = 0$

This takes one step for each column: $m = 3$ steps on elements (1, 1), (2, 2) and (3, 6)

$$\begin{aligned} \min f &= -6x_1 - 5x_3 - \frac{7}{3}x_4 + \frac{8}{3}x_5 + \frac{26}{3} \\ 3x_1 + 3x_3 + x_4 - 2x_5 &= 2 \\ 3x_1 + 3x_2 + x_4 + x_5 &= 8 \\ 6x_3 + x_4 - 5x_5 + 3x_6 &= 1 \\ x_1, \dots, x_6 &\geq 0 \end{aligned}$$

$-\frac{26}{3}$	-6	0	-5	$-\frac{7}{3}$	$\frac{8}{3}$	0
2	3	0	3	1	-2	0
8	3	3	0	1	1	0
1	0	0	6	1	-5	3

$$\begin{aligned} \min f &= x_3 - \frac{1}{3}x_4 - \frac{4}{3}x_5 + \frac{14}{3} \\ x_1 + x_3 + \frac{1}{3}x_4 - \frac{2}{3}x_5 &= \frac{2}{3} \\ -x_3 + x_5 &= 2 \\ 2x_3 + \frac{1}{3}x_4 - \frac{5}{3}x_5 + x_6 &= \frac{1}{3} \\ x_1, \dots, x_6 &\geq 0 \end{aligned}$$

$\frac{14}{3}$	0	0	1	$-\frac{1}{3}$	$-\frac{4}{3}$	0
$\frac{2}{3}$	1	0	1	$\frac{1}{3}$	$-\frac{2}{3}$	0
2	0	1	-1	0	1	0
$\frac{1}{3}$	0	0	2	$\frac{1}{3}$	$-\frac{5}{3}$	1

Graphical representation

$$\begin{aligned} \min f &= -x_1 - x_2 \\ 6x_1 + 4x_2 + x_3 &= 24 \\ 3x_1 - 2x_2 + x_4 &= 6 \\ x_1, x_2, x_3, x_4 &\geq 0 \end{aligned}$$

$-d$	x_1	x_2	x_3	x_4
0	-1	-1	0	0
24	6	4	1	0
6	3	-2	0	1

The basis is $(3, 4)$ and the basic solution $(0, 0, 24, 6)$

The basic reduced costs are zero, and the objective value is 0

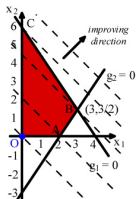
Perform *pivot* operations on $(1, 1)$ and $(2, 2)$ to obtain basis $(1, 2)$

$-d$	x_1	x_2	x_3	x_4
$\frac{9}{2}$	0	0	$\frac{5}{24}$	$-\frac{1}{12}$
3	1	0	$\frac{1}{12}$	$\frac{1}{6}$
$\frac{3}{2}$	0	1	$\frac{1}{8}$	$-\frac{1}{4}$

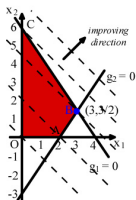
The basis is $(1, 2)$ and the basic solution is $(3, 3/2, 0, 0)$

The basic reduced costs are zero, and the objective value is $-\frac{9}{2}$

$$\begin{aligned} \min f &= -x_1 - x_2 \\ 6x_1 + 4x_2 &\leq 24 \\ 3x_1 - 2x_2 &\leq 6 \\ x_1, x_2 &\geq 0 \end{aligned}$$



The current vertex is $(0, 0)$



The current vertex is $(3, 3/2)$

Moving from a basis to another basis

When the problem is already in basic canonical form
a *pivot* operation on element a_{ij} represents a basis change

- the *pivot* column j enters the basis
- the basic column h with coefficient $a_{ih} = 1$ in the *pivot* row i exits

Therefore, the *pivot* element determines the resulting basis $\bar{B} = B \cup \{j\} \setminus \{h\}$

We will follow the evolution of the algorithm from three parallel points of view

- 1 algebraic representation: variables and equations
- 2 programming representation: *tableau*
- 3 graphical representation: possible only for 2 natural variables

A wrong choice of the *pivot* element

For example, we can move from basis (3, 4) to basis (2, 3):

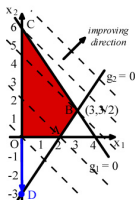
- x_2 gets into the basis: the *pivot* column is 2
- x_4 gets out of the basis: the *pivot* row is 2 (only nonzero coefficient)

the *pivot* element is $a_{22} = -1$

This choice of the *pivot* element is completely wrong:

- it produces an unfeasible solution:
if x_4 goes to zero, x_1 is unchanged and
 $3x_1 - 2x_2 + x_4 = 6$, then x_2 decreases
 $\bar{x}_2 = \bar{b}_2 := b_2/a_{22} = 6 / -2 = -3 < 0$
- it produced a worse solution:
if x_1 is unchanged, x_2 decreases and
 $f(x) = -x_1 - x_2$, then the cost increases
 $-\bar{d}_j := -d + (c_j/a_{ij}) b_i = 0 + (-1/-2) 6 = 3$

See also the graphic



The current vertex is
(0, -3)

How to choose the *pivot* element

At first, we introduce a fundamental assumption: **let the current basic solution be feasible** ($b_i \geq 0$ for all i) (we will see later how to obtain it)

The choice of the new basis, that is of the *pivot* element, must guarantee that

① the new basic variable remains nonnegative

- algebraic perspective: if the old basic variable decreases to zero, the new one must increase
- *tableau* perspective: $\bar{b}_i := \frac{b_i}{a_{ij}} \geq 0$
- graphical perspective: we must move towards the feasible region

Therefore, **choose a positive pivot element: $a_{ij} > 0$**

② the objective function improves

- algebraic perspective: when the new basic variable increases, the total cost must decrease
- *tableau* perspective: $-\bar{d} = -d - \frac{c_j}{a_{ij}} b_i > -d$
- graphical perspective: we must move consistently with the most improving direction

Therefore, **choose a pivot column with negative reduced cost: $c_j < 0$**

How to choose the *pivot* element

The choice of the new basis, that is of the *pivot* element, must guarantee that

③ **the other basic variables remain feasible** (one decreases to zero)

- algebraic perspective: if x_j increases by ϵ , and the nonbasic variables keep zero, each old basic variable changes by $-\epsilon a_{hj}$ ($h = 1, \dots, m$); they keep feasible as long as

$$x_{B_h}^- = b_h - \epsilon a_{hj} \geq 0 \Leftrightarrow \epsilon a_{hj} \leq b_h \text{ for all } h \neq i$$

- *tableau* perspective:

$$\bar{b}_h := b_h - \frac{a_{hj}}{a_{ij}} b_i \geq 0 \Leftrightarrow \begin{cases} \frac{b_i}{a_{ij}} \leq \frac{b_h}{a_{hj}} & \text{for } a_{hj} > 0 \\ \frac{b_i}{a_{ij}} \geq 0 & \text{for } a_{hj} = 0 \text{ (trivial)} \\ \frac{b_i}{a_{ij}} \geq \frac{b_h}{a_{hj}} & \text{for } a_{hj} < 0 \text{ (trivial)} \end{cases}$$

- graphical perspective:

- if $a_{hj} < 0$, we move away from the separating hyperplane $a_h^T x = b_h$
- if $a_{hj} = 0$, we move parallel to the separating hyperplane $a_h^T x = b_h$
- if $a_{hj} > 0$, we move closer to the separating hyperplane $a_h^T x = b_h$

Therefore, **choose the *pivot* row with minimum b_h/a_{hj} among those with $a_{hj} > 0$**

$$i := \arg \min_{i: a_{ij} > 0} \frac{b_h}{a_{hj}}$$

Example (1)

$$\begin{aligned}\min f &= -x_1 - x_2 \\ 6x_1 + 4x_2 + x_3 &= 24 \\ 3x_1 - 2x_2 + x_4 &= 6 \\ x_1, x_2, x_3, x_4 &\geq 0\end{aligned}$$

$-d$	x_1	x_2	x_3	x_4
0	-1	-1	0	0
24	6	4	1	0
6	3	-2	0	1

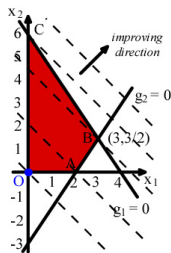
The basis is (3, 4)

The basic solution (0, 0, 24, 6)

The objective value is 0

The basic reduced costs are zero

$$\begin{aligned}\min f &= -x_1 - x_2 \\ 6x_1 + 4x_2 &\leq 24 \\ 3x_1 - 2x_2 &\leq 6 \\ x_1, x_2 &\geq 0\end{aligned}$$



The current vertex is (0, 0)

Two nonbasic variables have the same negative reduced cost: $c_1 = c_2 = -1$

Variable x_1 enters the basis (x_2 would be better, but the selection rule is heuristic)

The ratio b_i/a_{ij} is minimum for $i = 2$: variable x_4 leaves the basis; **pivot element a_{21}**

Example (2)

$$\begin{aligned}\min f &= -x_1 - x_2 \\ 6x_1 + 4x_2 + x_3 &= 24 \\ 3x_1 - 2x_2 + x_4 &= 6 \\ x_1, x_2, x_3, x_4 &\geq 0\end{aligned}$$

$-d$	x_1	x_2	x_3	x_4
2	0	$-\frac{5}{3}$	0	$\frac{1}{3}$
12	0	8	1	-2
2	1	$-\frac{2}{3}$	0	$\frac{1}{3}$

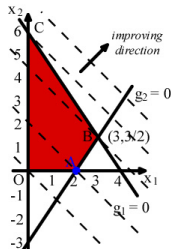
The basis is (1, 3)

The basic solution (2, 0, 12, 0)

The objective value is -2

The basic reduced costs are zero

$$\begin{aligned}\min f &= -x_1 - x_2 \\ 6x_1 + 4x_2 &\leq 24 \\ 3x_1 - 2x_2 &\leq 6 \\ x_1, x_2 &\geq 0\end{aligned}$$



The current vertex is (2, 0)

Only one nonbasic variable has a negative reduced cost: $c_2 = -5/3$

Variable x_2 enters the basis

The ratio b_i/a_{ij} is positive only for row $i = 1$: variable x_3 leaves the basis;

pivot element a_{12}

Example (3)

$$\begin{aligned} \min f &= -x_1 - x_2 \\ 6x_1 + 4x_2 + x_3 &= 24 \\ 3x_1 - 2x_2 + x_4 &= 6 \\ x_1, x_2, x_3, x_4 &\geq 0 \end{aligned}$$

$-d$	x_1	x_2	x_3	x_4
9	0	0	5	1
2			24	12
3	0	1	1	1
2			8	4
3	1	0	1	1
			12	6

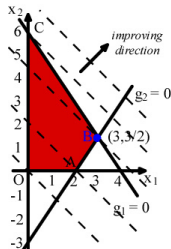
The basis is (1, 2)

The basic solution is $(3, 3/2, 0, 0)$

The objective value is $-9/2$

The basic reduced costs are zero

$$\begin{aligned} \min f &= -x_1 - x_2 \\ 6x_1 + 4x_2 &\leq 24 \\ 3x_1 - 2x_2 &\leq 6 \\ x_1, x_2 &\geq 0 \end{aligned}$$



The current vertex is $(3, 3/2)$

Only one nonbasic variable has a negative reduced cost: $c_4 = -1/12$

Variable x_4 enters the basis

The ratio b_i/a_{ij} is positive only for row $i = 2$: variable x_1 leaves the basis;

pivot element a_{24}

Example (4)

$$\begin{aligned} \min f &= -x_1 - x_2 \\ 6x_1 + 4x_2 + x_3 &= 24 \\ 3x_1 - 2x_2 + x_4 &= 6 \\ x_1, x_2, x_3, x_4 &\geq 0 \end{aligned}$$

$-d$	x_1	x_2	x_3	x_4
6	$\frac{1}{2}$	0	$\frac{1}{4}$	0
6	$\frac{3}{2}$	1	$\frac{1}{4}$	0
18	6	0	$\frac{1}{2}$	1

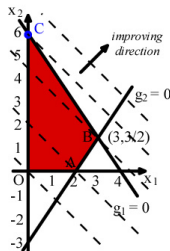
The basis is (2, 4)

The basic solution (0, 6, 0, 18)

The objective value is -6

The basic reduced costs are zero

$$\begin{aligned} \min f &= -x_1 - x_2 \\ 6x_1 + 4x_2 &\leq 24 \\ 3x_1 - 2x_2 &\leq 6 \\ x_1, x_2 &\geq 0 \end{aligned}$$



The current vertex is (0, 6)

All reduced costs are nonnegative: **the current solution is optimal**