## Foundations of Operations Research

## Master of Science in Computer Engineering

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Lesson 1: Introduction to Operations Research

## Operations Research

It is the science of decision-making
It helps people taking better decisions in complex situations, exploiting

- mathematical models
- quantitative methods
to determine the results of each possible decision and
to select the best decision (or at least a satisfactory one)
It is a branch of applied mathematics combining
- economics (to model the decision process)
- knowledge from specific application fields such as industrial engineering, marketing, physics, chemistry, biology, etc. . . (to describe correctly the problems)
- computer science (to solve the problems efficiently)
- business (to turn the suggested decisions into practice)


## Decision-making problems

They are characterized by two relevant features
(1) the number of solutions:

- many classical mathematical problems have a single solution:

Apples cost 1.5 euro/kg.
You buy 3 kg. of apples
How much do you spend?

- decision-making problems have several feasible solutions and some solutions are better than others:

Go to the market
Buy enough food for the whole week
Spend as little as possible
(2) the impact on practical actions:

- the solution to a classical mathematical problem is
- an answer to an abstract question
- an information on what has happened or will happen
- the optimal solution to a decision-making problem is a guide about how to act in a practical situation


## Example: task assignment

- $m=3$ tasks have to be performed: $T_{1}, \ldots, T_{m}$
- $m=3$ machines are available to perform them: $M_{1}, \ldots, M_{m}$
- each task must be assigned to exactly one machine
- each machine can perform exactly one task
- $c_{i j}$ is the cost required to perform task $T_{i}$ on machine $M_{j}$ $(i, j=1, \ldots, m)$

|  | $M_{1}$ | $M_{2}$ | $M_{3}$ |
| :---: | :---: | :---: | :---: |
| $T_{1}$ | 2 | 6 | 3 |
| $T_{2}$ | 8 | 4 | 9 |
| $T_{3}$ | 5 | 7 | 8 |

Decide which machine will perform each task so as to minimize the total cost

How many solutions are there?

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## Example: network design

- $n=5$ towns (nodes) must be linked to each other (directly or not)
- $m=9$ potential links (edges) exist between towns $i$ and $j$
- each potential link $(i, j)$ has a cost $c_{i j}$


Select a subset of edges guaranteeing that each pair of nodes are linked so as to minimize the total cost

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At most $2^{m}$, that is the number of subsets of links (not all subsets guarantee the connectivity)

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## Shortest path

- a digital map of a road network
- an origin location
- a target location


Determine a route from the origin to the target
so as to minimize the travel time (or distance, or cost)

## Social networks

A social network

- a set of users
- a set of friendship relations between pairs of users


Find a subset of users who are all pairwise related so as to maximize their number

## Example: mail delivery

- a digital map of a road network
- a starting and arrival point
- a set of streets along which mail must be delivered


Find a close path along all the required streets so as to minimize its total length

## Examples

- Personnel scheduling: determine the weekly shifts for workers covering all daily tasks and respecting labor regulations so as to minimize the cost of the personnel
- Service counters:
determine how many counters to open at each time of the day respecting a limit on the average waiting time of the customers so as to minimize the cost
- Laptop: choose the components of a laptop so as to optimize its total price, weight and performance


## A brief historical sketch: the roots

Combinatorial
Optimization


Leonhard Euler (1707-1783)

Linear
Programming


Jean Baptiste Joseph Fourier (1768-1830)

Integer
Programming


Diophantus (200/214-284/298)
(1) 1735: Euler solves the problem of the seven bridges of Königsberg
(2) 1826: Fourier solves 3-dimensional Linear Programming problems
(3) III century: Diophantus searches for integer solutions to equations

## A brief historical sketch: the beginnings

World War II gave the real start to Operations Research, summoning scientists to research on the most efficient conduct of military operations

During the Siege of Leningrad, L.V. Kantorovich (19121986) was in charge of the Road of Life, an ice road across the frozen Lake Ladoga, which was the only access to the town during winter.

He calculated the optimal distance between cars on ice, depending on ice thickness and air temperature, and personally walked on the ice between the cars to ensure they did not sink.


## A brief historical sketch: the developments

Combinatorial Optimization:

- 1926: O. Borúvka finds how to connect all nodes in a network at minimum cost
- 1947: E. Dijkstra finds the shortest path between two places

Linear Programming

- 1939: L.V. Kantorovitch lays the foundations of Linear Programming (Nobel Prize in 1975)
- 1947: G. Dantzig invents the simplex algorithm for linear programs

Integer Programming

- 1958: R. Gomory proposes the cutting plane method
- 1965: E. Balas proposes the branch-and-bound method


## A brief historical sketch: business applications

After the war, the substantial increase in the size of companies and organizations gave rise to more complex decision-making problems

OR was widely applied to business, industry, and society

- fast progress in methodologies
- diffusion of computing power (hardware and software)

OR methodologies improve the use of scarse resources with a significant impact not only for large companies and organizations

Operations Research $=$ Management Science
Rapidly evolving contexts, with high levels of complexity and uncertainty pose a severe challenge to Operations Research

The huge amount of data available (Big Data) with modern information systems opens new avenues (Business Analytics)

## A brief historical sketch: business applications

## OPERATIONS RESEARCH: THE SCIENCE OF BETTER

tIME-STARVED EXECUTIVES ARE MAKING BOLDER DECISIONS WITH LESS RISK aND BETTER OUTCOMES. THEIR SECRET: OPERATIONS RESEARCH.

| year | company | sector | results |
| :---: | :---: | :---: | :---: |
| 1990 | Taco Bell <br> (fast food) | personnel scheduling | $7.6 \mathrm{M} \mathrm{\$}$ annual <br> savings |
| 1992 | American Arlines | design fare structure, overbooking <br> and flights coordination | $+500 \mathrm{M} \$$ |
| 1992 | Harris Corp. <br> (semiconductors) | production planning | $50 \% \Rightarrow 95 \%$ orders <br> on time |
| 1995 | GM - car rental | use of car park | $+50 \mathrm{M} \$$ per year <br> avoided bankruptcy |
| 1996 | HP - printers | modify production line | doubled <br> production |
| 1997 | Bosques Arauco | harvesting logistics <br> and transport | $5 M \$$ annual <br> savings |
| 1999 | IBM | supply chain re-engineering <br> savings | $750 \mathrm{M} \$$ annual <br> sang |

## Not just money-making

Redesigning Milan's Emergency Service 118


